

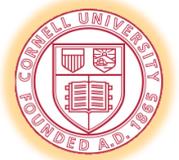
A Compiler and Run-time System for Network Programming Languages

Christopher Monsanto, Princeton

Nate Foster, Cornell

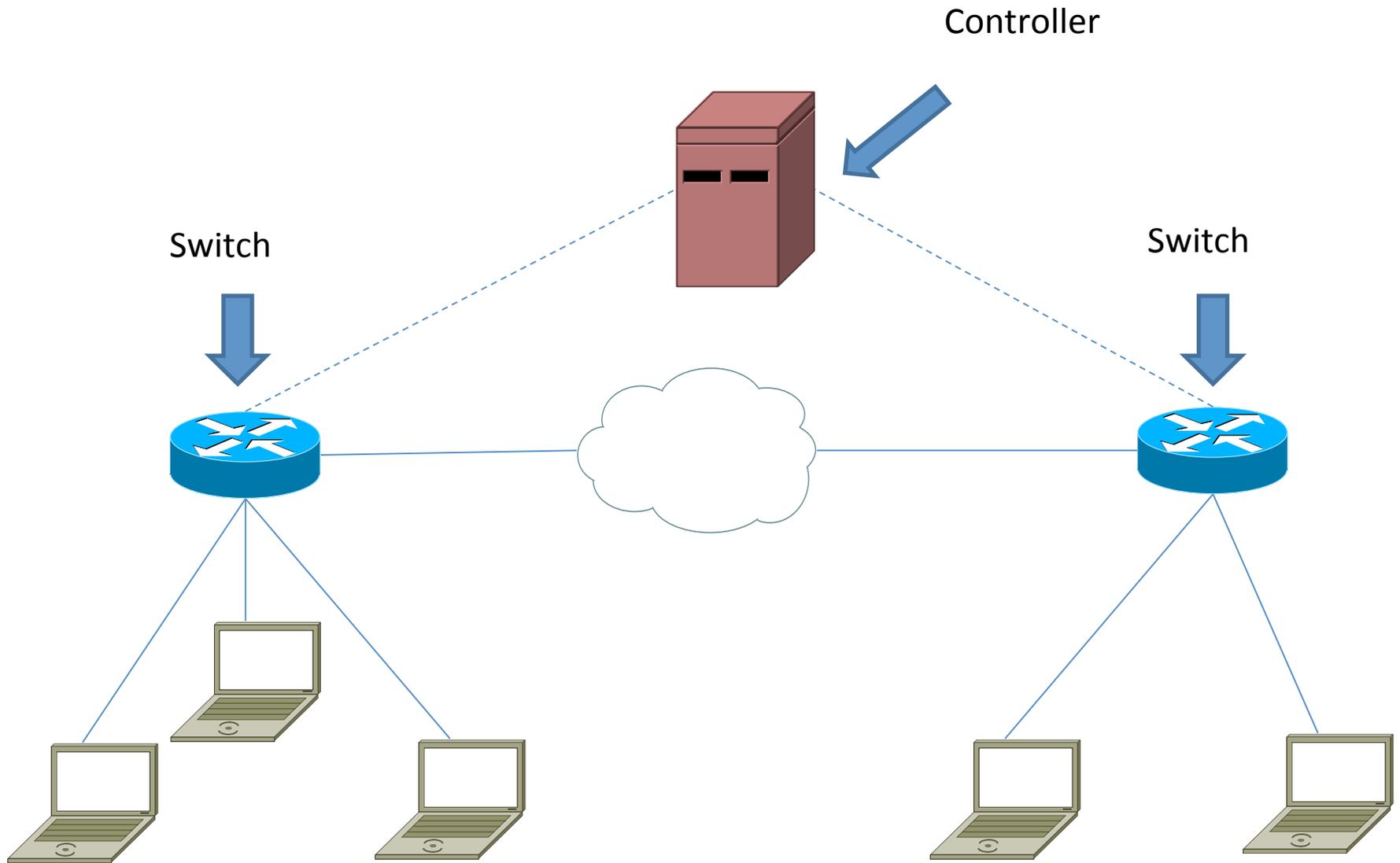
Rob Harrison, West Point

David Walker, Princeton



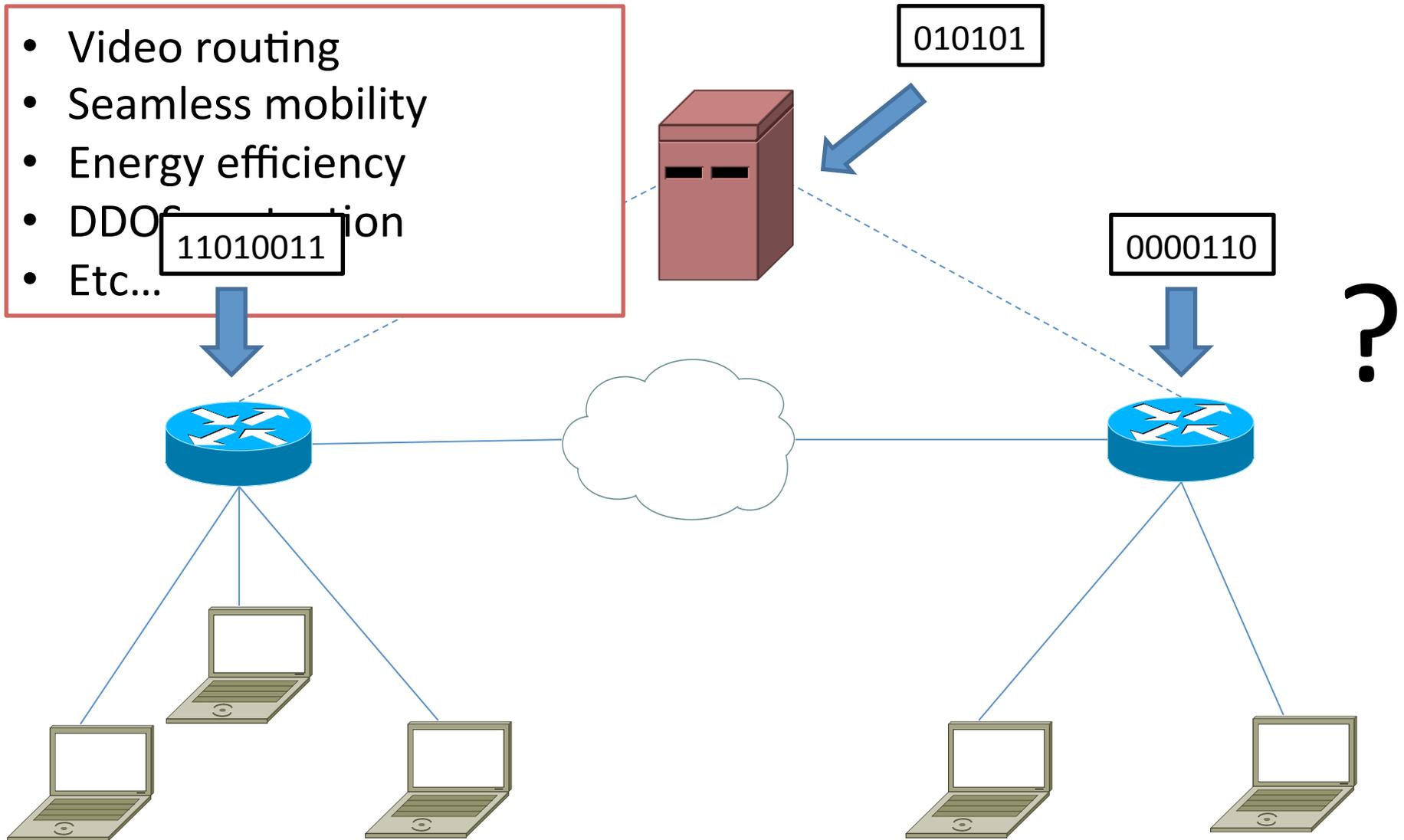
frenetic >>

Software-Defined Networks



Software-Defined Networks

- Video routing
- Seamless mobility
- Energy efficiency
- DDoS mitigation
- Etc...



Software Defined Networks: Switches

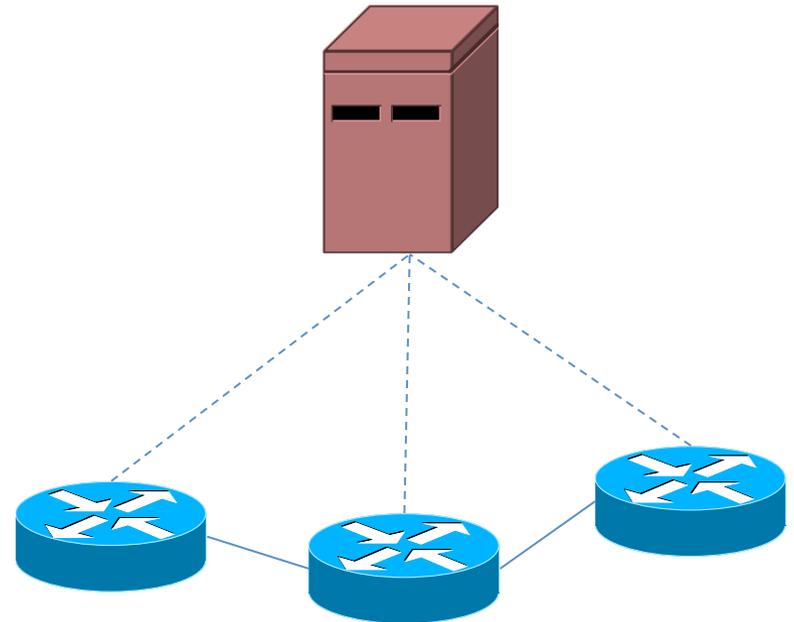
- Processes packets according to classifier
 - Limited pattern matching capabilities
 - Actions: forwarding, dropping, sending to controller
- As fast as it gets

| Dest. IP | Dest. Port | Action |
|-----------|------------|------------|
| 191.*.*.* | 80 | Forward A |
| 191.*.*.* | * | Drop |
| 191.1.1.0 | 22 | Forward B |
| * | * | Controller |



SDN: Controllers

- Capable of arbitrary computation
- Orders of magnitude slower
- Installs and uninstalls rules from switch classifiers



Goals

1. A simple, declarative language...
 - That is, we specify the functionality, not the concrete rules that go on the switches, or the explicit install and uninstall commands the controller must issue
2. ... that is mathematically guaranteed to be correct and efficient.

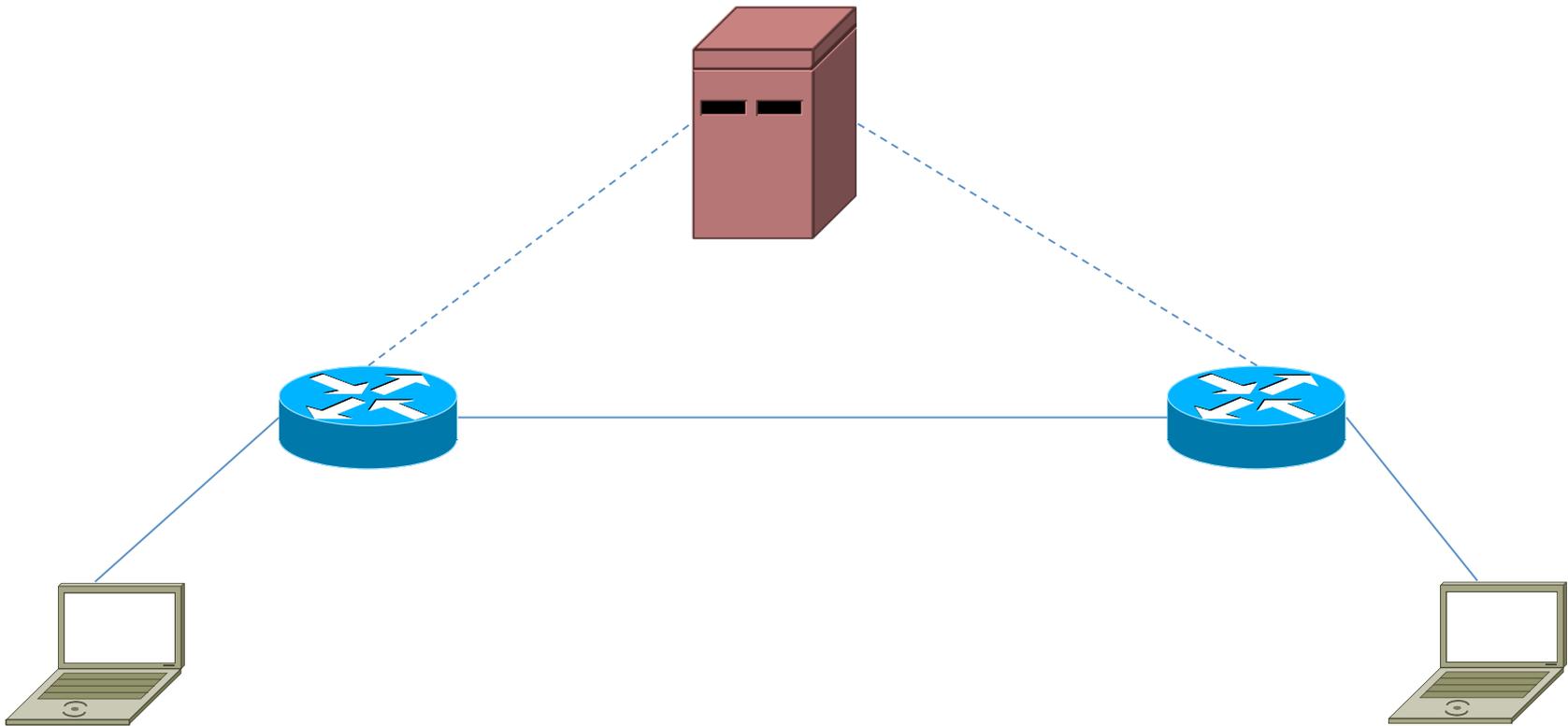
Goals

1. A simple, declarative language...
 - That is, we specify the functionality, not the concrete rules that go on the switches, or the explicit install and uninstall commands the controller must issue
2. ... that is mathematically guaranteed to be correct and *efficient*.

???

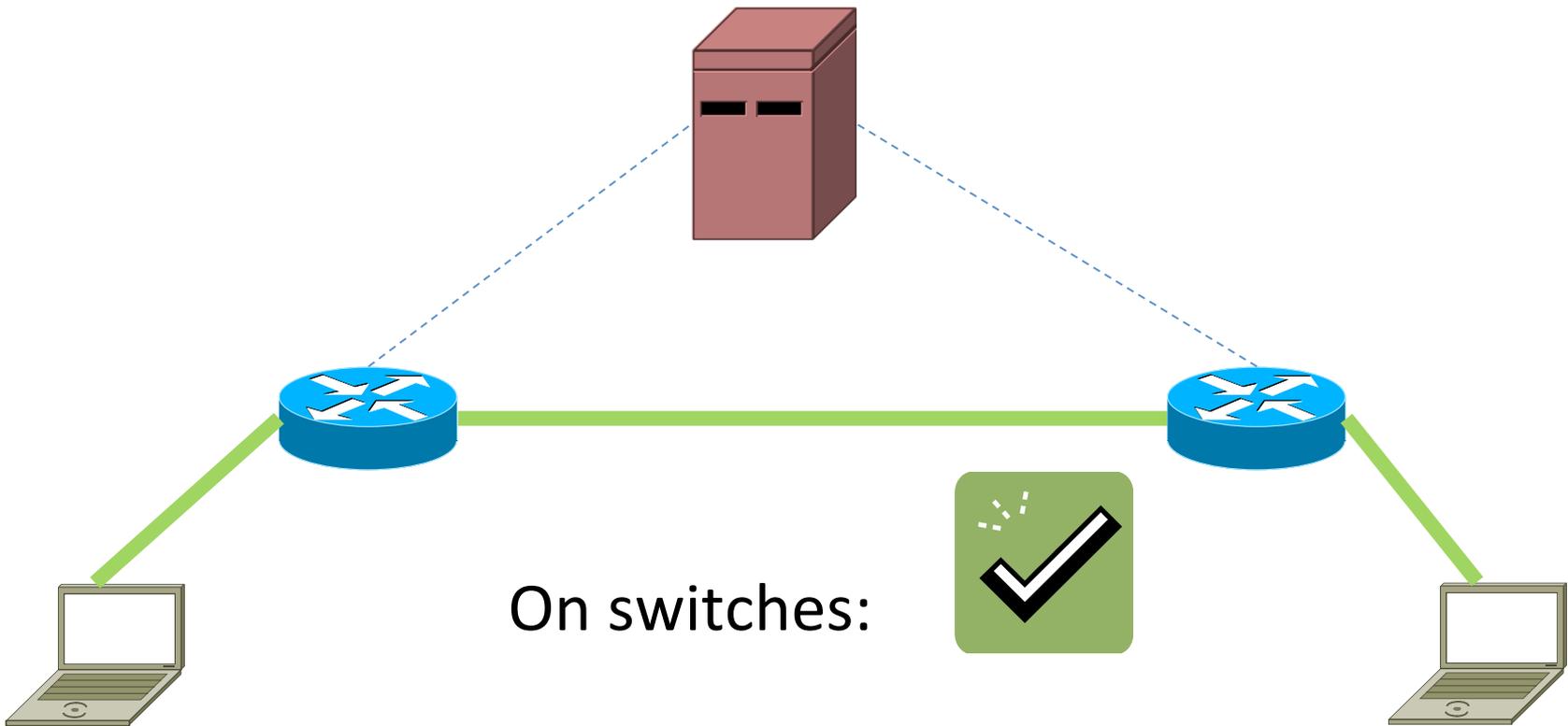
Goals

- ... guaranteed to be *efficient*



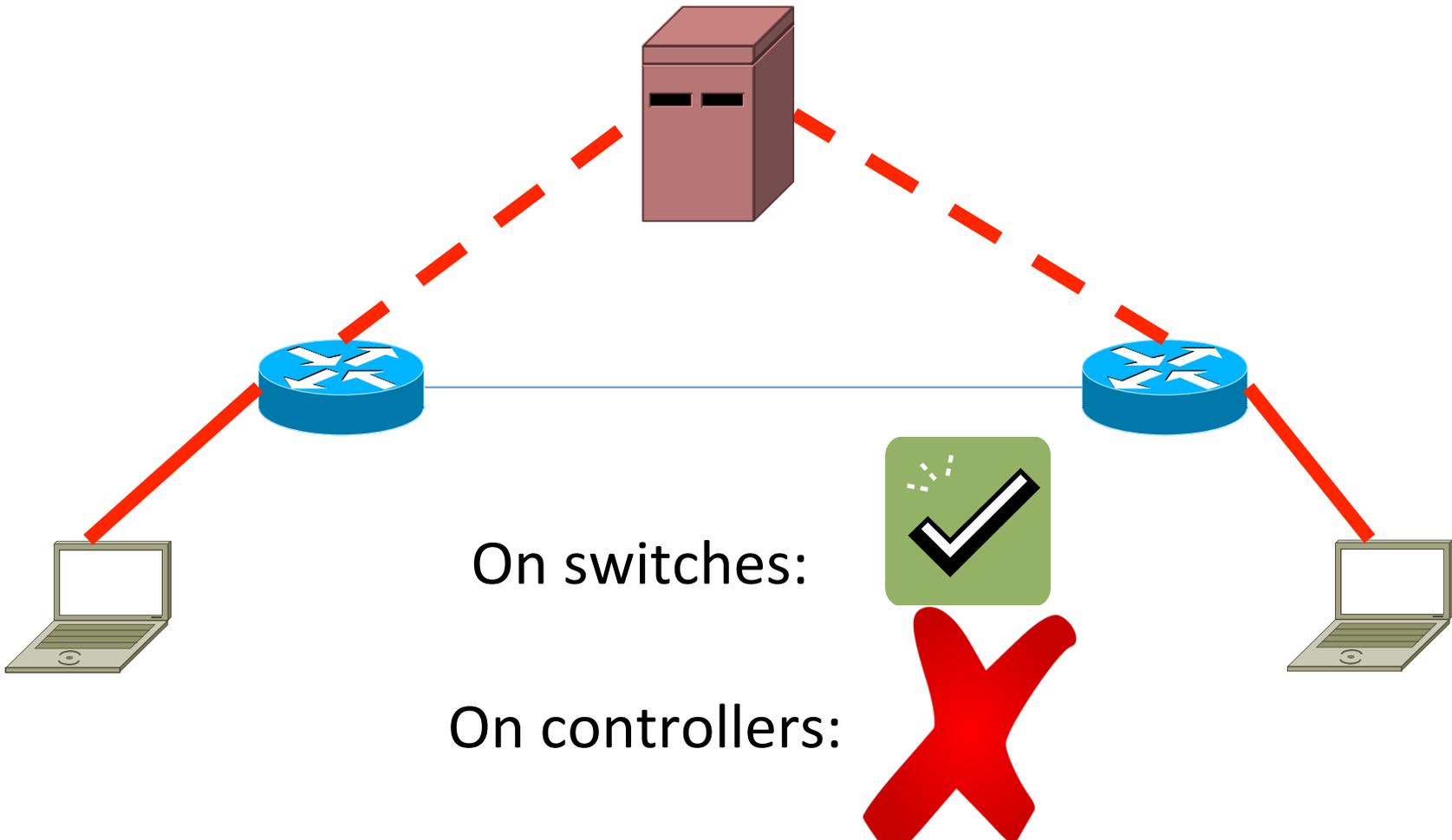
Goals

- ... guaranteed to be *efficient*



Goals

- ... guaranteed to be *efficient*

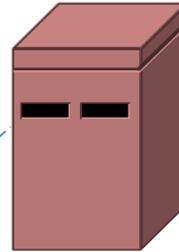


Obstacles Faced By Programmers

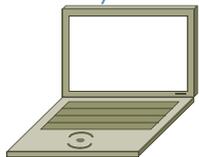
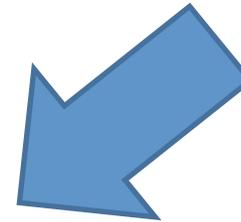


Low level interface

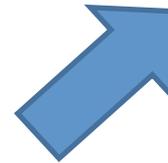
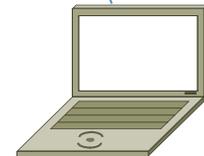
| Dest. IP | Dest. Port | Action |
|-----------|------------|------------|
| 191.*.*.* | 80 | Forward(A) |
| 191.*.*.* | * | Drop |
| * | * | Controller |



Synthesize communication protocol



Differing switch capabilities

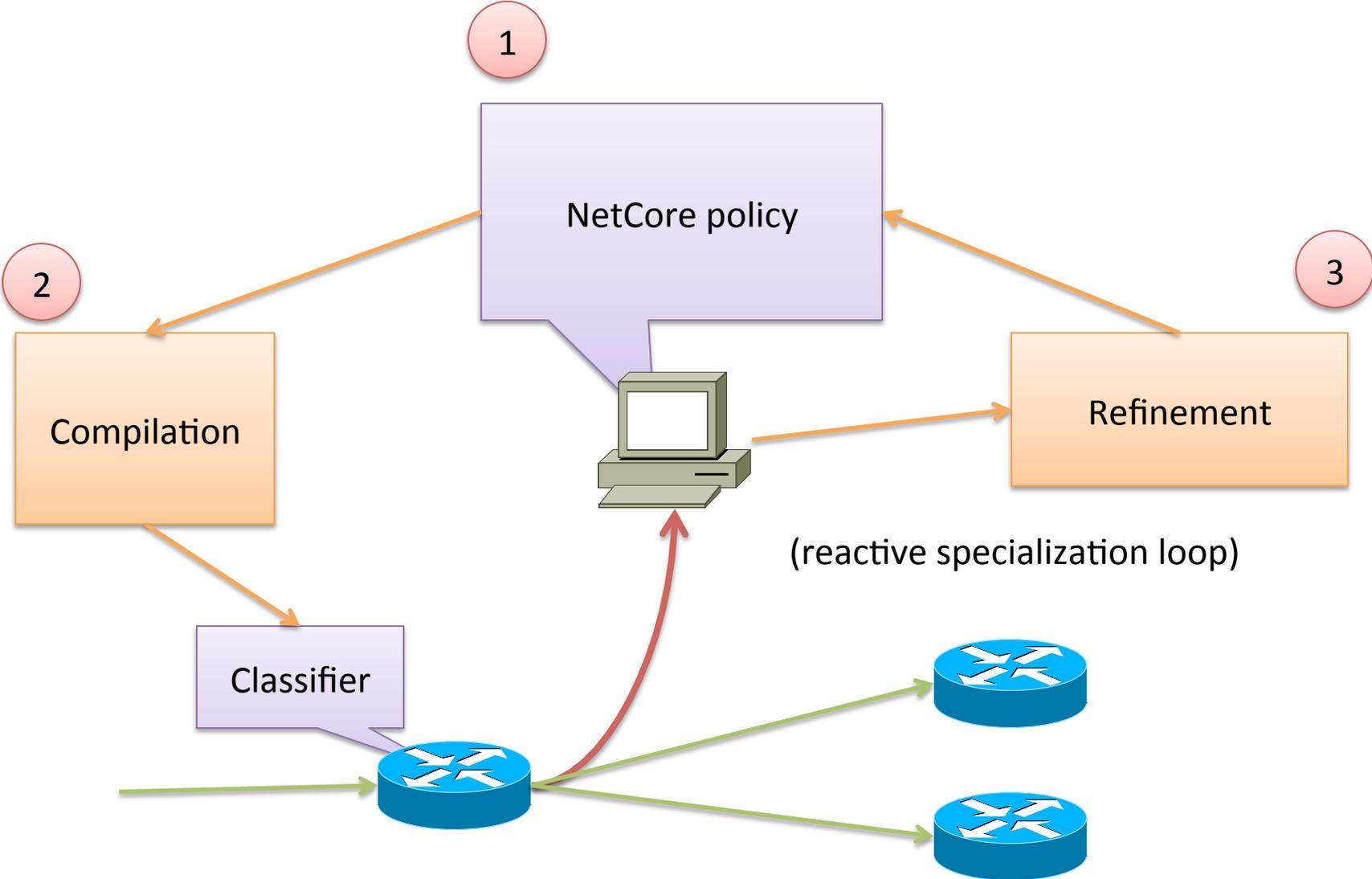


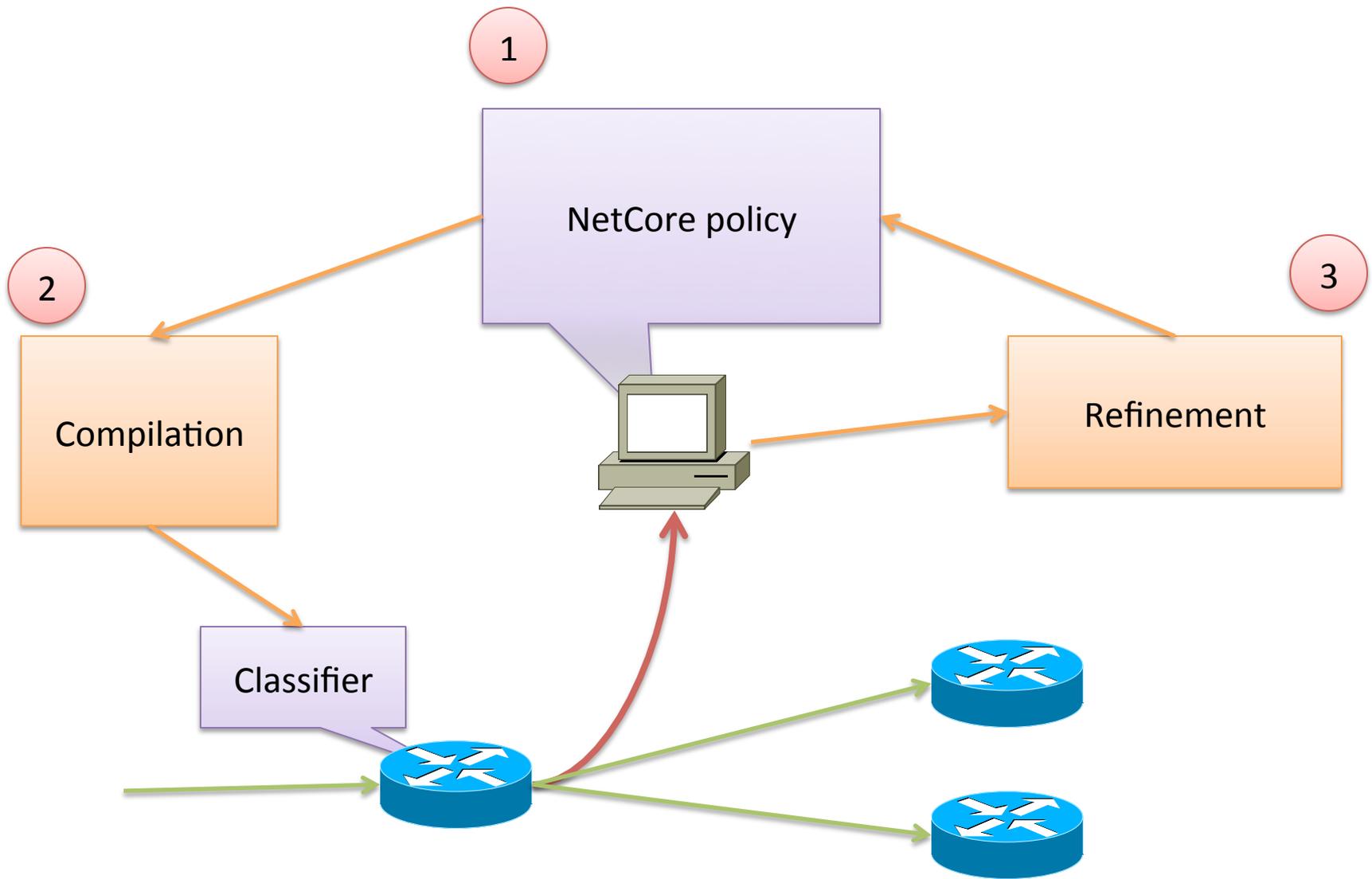
Our Real Enemy: Complexity

- Managing these rules is complicated
- Tendency to fall back to using the simplest kind of rule: microflows (exact match rules)
- Very inefficient: many packets go to the controller!



New Frenetic Run-time Architecture





NETCORE

NetCore Grammar

Predicates

- $e ::= h:w$
 - | switch s
 - | inspect f
 - | $e \downarrow 1 \cup e \downarrow 2$
 - | $e \downarrow 1 \setminus e \downarrow 2$
 - | $e \downarrow 1 \cap e \downarrow 2$
 - | $\neg e$

Policies

- $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action
 - | $\tau \downarrow 1 \cup \tau \downarrow 2$

Headers h , wildcards w

- primitive match
- match packets at switch
- arbitrary function
- set union
- set difference
- set intersection
- set negation

policy union

NetCore Grammar

Predicates

- $e ::= h:w$
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switch s

inspect f

$e \downarrow 1 \cup e \downarrow 2$

$e \downarrow 1 \setminus e \downarrow 2$

$e \downarrow 1 \cap e \downarrow 2$

$\neg e$

Policies

• $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action

$\tau \downarrow 1 \cup \tau \downarrow 2$

Headers h , wildcards w

primitive match

DestIP: 10.*.1.*

arbitrary function

set union

set difference

set intersection

set negation

policy union

NetCore Grammar

Predicates

Headers h , wildcards w

- $e ::= h:w$
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arbitrary function
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NetCore Grammar

Predicates

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Policies

- $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action
- | $\tau \downarrow 1 \cup \tau \downarrow 2$
policy union

NetCore Grammar

Predicates

Headers h , wildcards w

• $e ::= h:w$

primitive match

| switch s

match packets at switch

| inspect f

arbitrary function

```
badWebsite p = addr `elem` blacklist
```

where

```
headers = parseHTTP p
```

```
addr = headers ! "address"
```

Policies

• $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action

| $\tau \downarrow 1 \cup \tau \downarrow 2$

policy union

NetCore Grammar

Predicates

Headers h , wildcards w

• $e ::= h:w$

primitive match

| switch s

match packets at switch

| inspect f

arbitrary function

```
badWebsite p = addr `elem` blacklist
```

where

```
headers = parseHTTP p
```

```
addr = headers ! "address"
```

Policies

```
inspect (not . badWebsite)
```

• $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action

| $\tau \downarrow 1 \cup \tau \downarrow 2$

policy union

NetCore Grammar

Predicates

- $e ::= h:w$ Headers h , wildcards w
primitive match
 - | switch s match packets at switch
 - | inspect f arbitrary function
 - | $e \downarrow 1 \cup e \downarrow 2$ set union
 - | $e \downarrow 1 \setminus e \downarrow 2$ set difference
 - | $e \downarrow 1 \cap e \downarrow 2$ set intersection
 - | $\neg e$ set negation
- Policies

- $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action
- | $\tau \downarrow 1 \cup \tau \downarrow 2$ policy union

NetCore Grammar

Predicates

Headers h , wildcards w

- $e ::= h:w$ primitive match
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- | $e \downarrow 1 \cup e \downarrow 2$ set union
- | $e \downarrow 1 \setminus e \downarrow 2$ set difference
- | $e \downarrow 1 \cap e \downarrow 2$ set intersection

Policies

```
inspect (not . badWebSite) ∪  
DestIP: 10.10.50.20
```

- $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action
- | $\tau \downarrow 1 \cup \tau \downarrow 2$ policy union

NetCore Grammar

Predicates

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Policies

- $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action
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- | $\neg e$

Headers h , wildcards w

primitive match

match packets at switch

arbitrary function

set union

set difference

set intersection

set negation

Policies

- $\tau ::=$ inspect (not . badWebSite) \rightarrow {A}
- | $\tau \downarrow 1 \cup \tau \downarrow 2$

policy union

NetCore Grammar

Predicates

- $e ::= h:w$
- | switch s
- | inspect f
- | $e \downarrow 1 \cup e \downarrow 2$
- | $e \downarrow 1 \setminus e \downarrow 2$
- | $e \downarrow 1 \cap e \downarrow 2$
- | $\neg e$

Policies

- $\tau ::= e \rightarrow \{s \downarrow 1, \dots, s \downarrow n\}$ action
- | $\tau \downarrow 1 \cup \tau \downarrow 2$

Headers h , wildcards w

primitive match

match packets at switch

arbitrary function

set union

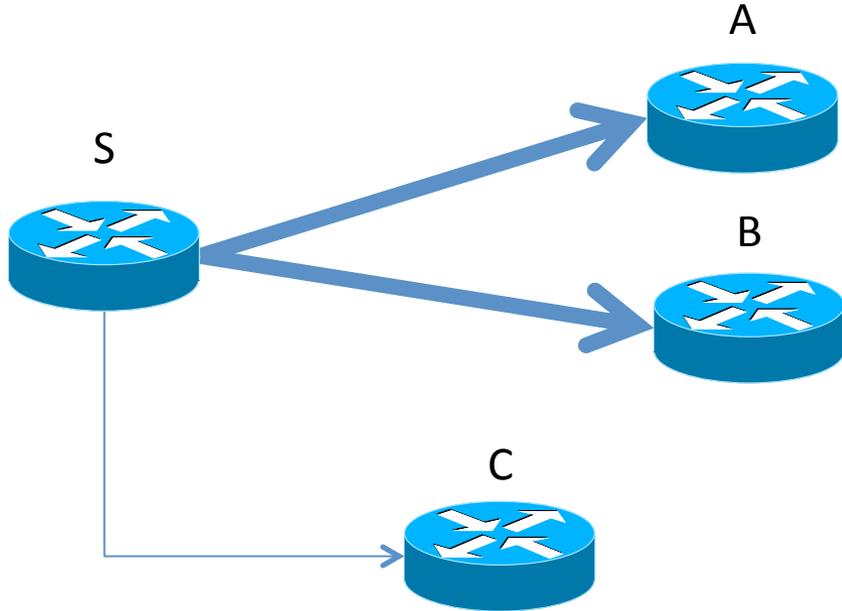
set difference

set intersection

set negation

policy union

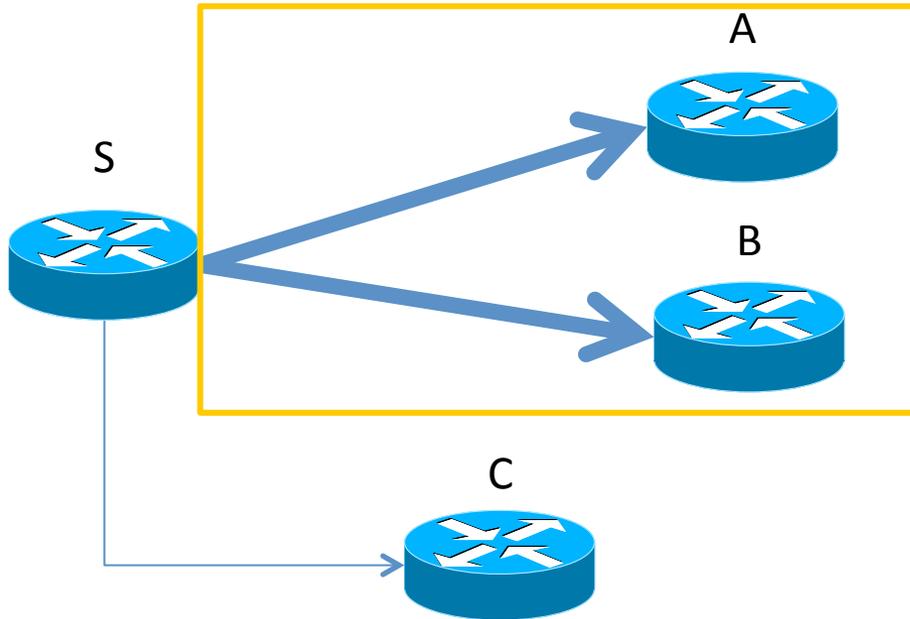
NetCore Example



- Load balanced, fast path
- Slow path

DestIP: *0 → {A} ∪ DestIP: *1 → {B}

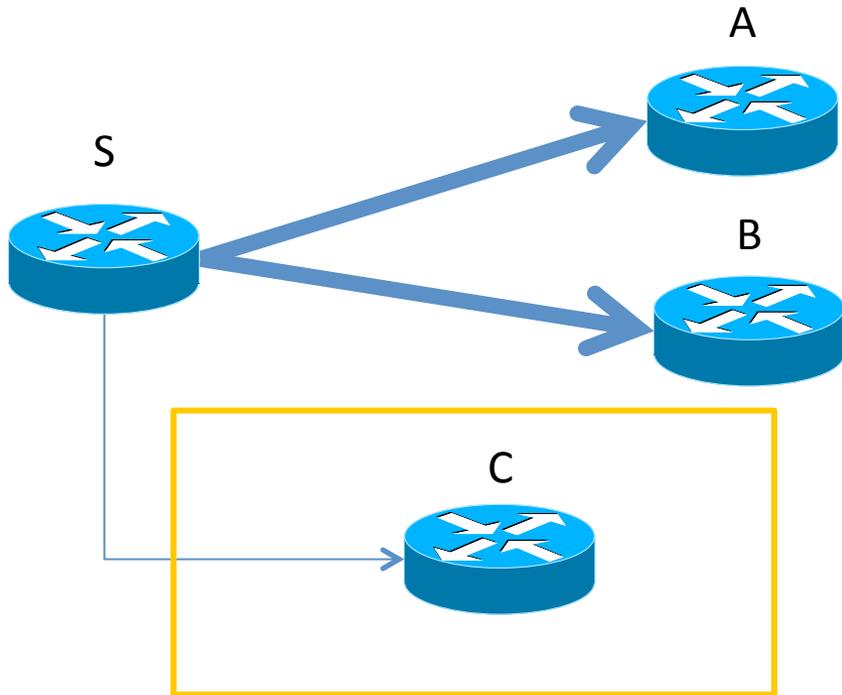
NetCore Example



- Load balanced, fast path
 - Internal traffic (90.80.*.*)
 - SSH traffic
- Slow path

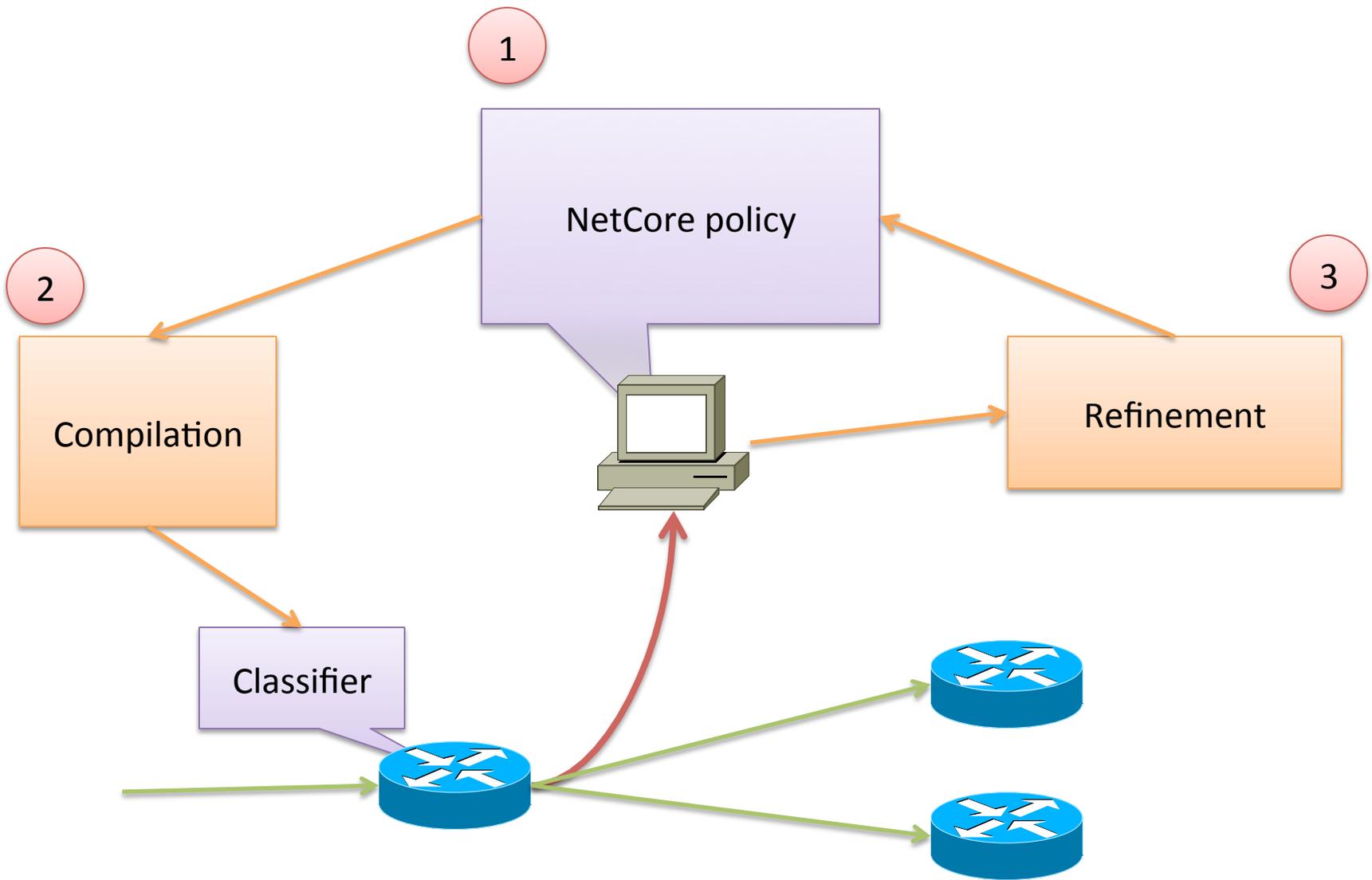
$(\text{DestIP: } 90.80.*.* \cup \text{DestPort: } 22) \cap$
 $(\text{DestIP: } *0 \rightarrow \{A\} \cup \text{DestIP: } *1 \rightarrow \{B\})$

NetCore Example



- Load balanced, fast path
 - Internal traffic (90.80.*.*)
 - SSH traffic
- Slow path
 - External traffic

```
((DestIP:90.80.*.* ∪ DestPort: 22) ∩  
  (DestIP: *0 → {A} ∪ DestIP: *1 → {B})) ∪  
  (¬DestIP:90.80.*.*) → {C}
```



COMPILATION

Compilation

P1 = DestIP:90.80.*.* → {A}

Compiles to:

| Dest. IP | Dest. Port | Action |
|-----------|------------|-----------|
| 90.80.*.* | * | Forward A |

Compilation: Approximations

Non prefix match!



P2 = DestIP:90.*.70.60 → {B}

Compiles to:

| Dest. IP | Dest. Port | Action |
|----------|------------|------------|
| 90.*.*.* | * | Controller |

Overapproximation



Compilation

Run programs in parallel:

$$P3 = P1 \cup P2$$

Compiles to:

| Dest. IP | Dest. Port | Action |
|-------------|------------|--------------|
| 90.80.70.60 | * | Forward A, B |
| 90.80.*.* | * | Forward A |
| 90.*.*.* | * | Controller |

Compilation

Run programs in parallel:

$$P3 = P1 \cup P2$$

Subtle interaction: where did B come from?



Compiles to:

| Dest. IP | Dest. Port | Action |
|-------------|------------|--------------|
| 90.80.70.60 | * | Forward A, B |
| 90.80.*.* | * | Forward A |
| 90.*.*.* | * | Controller |

Compilation: Exponential Growth

DestIP: 1.*.*.* → {A}

| Dest. IP | Dest. Port | Action |
|----------|------------|-----------|
| 1.*.*.* | * | Forward A |

Compilation: Exponential Growth

DestIP: 1.*.*.* \rightarrow {A} \cup
DestIP: *.2.*.* \rightarrow {B}

| Dest. IP | Dest. Port | Action |
|----------|------------|--------------|
| 1.2.*.* | * | Forward A, B |
| *.2.*.* | * | Forward B |
| 1.*.*.* | * | Forward A |

Compilation: Exponential Growth

DestIP: 1.*.*.* $\rightarrow \{A\} \cup$

DestIP: *.2.*.* $\rightarrow \{B\} \cup$

DestIP: *.*.3.* $\rightarrow \{C\}$

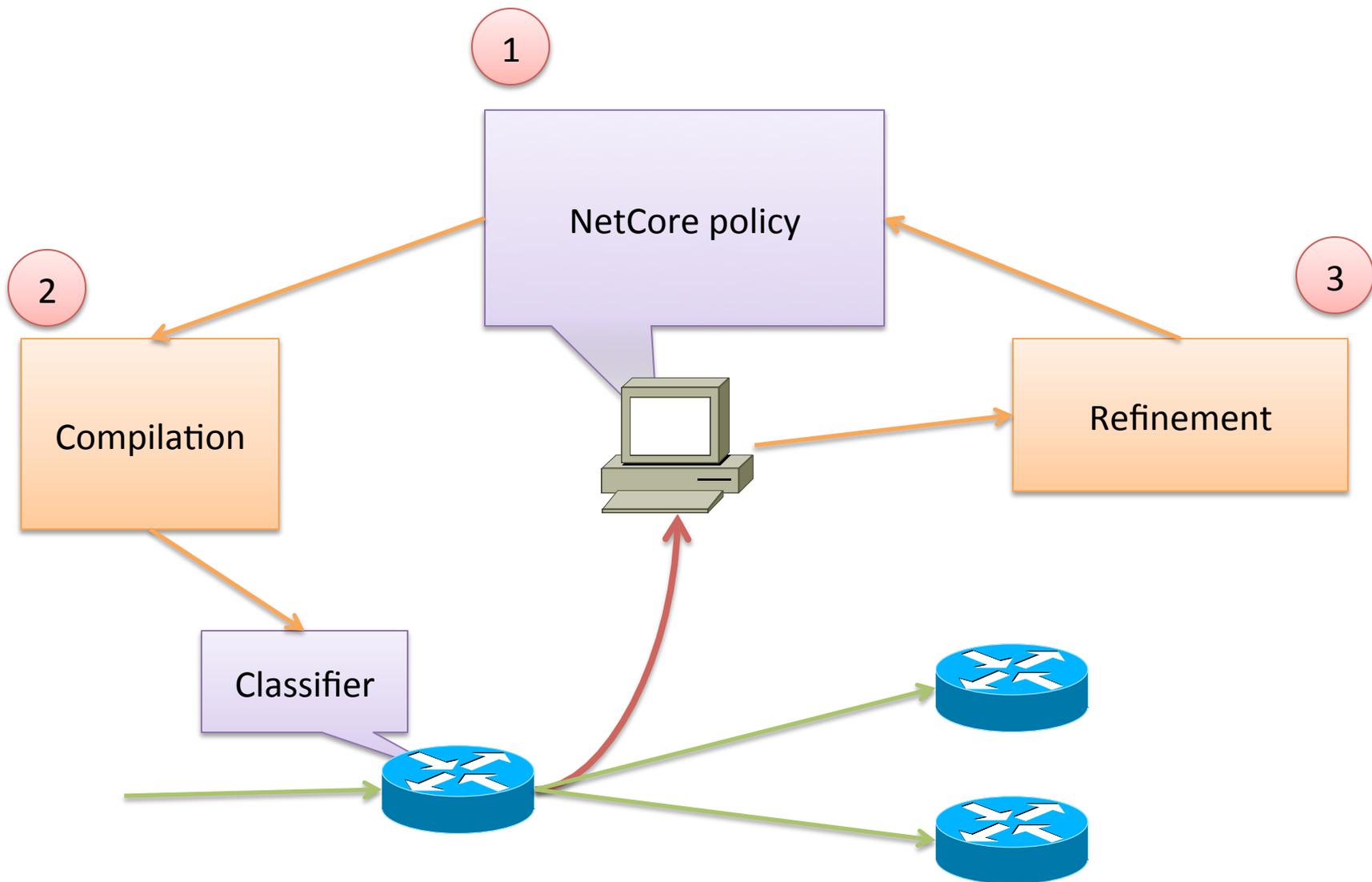
| Dest. IP | Dest. Port | Action |
|----------|------------|-----------------|
| 1.2.3.* | * | Forward A, B, C |
| *.2.3.* | * | Forward B, C |
| 1.*.3.* | * | Forward A, C |
| *.*.3.* | * | Forward C |
| 1.2.*.* | * | Forward A, B |
| *.2.*.* | * | Forward B |
| 1.*.*.* | * | Forward A |

Com... .. with

| Dest. IP | Dest. Port | Action |
|----------|------------|--------------------|
| 1.2.3.4 | * | Forward A, B, C, D |
| *.2.3.4 | * | Forward B, C, D |
| 1.*.3.4 | * | Forward A, C, D |
| *.*.3.4 | * | Forward C, D |
| 1.2.*.4 | * | Forward A, B, D |
| *.2.*.4 | * | Forward B, D |
| 1.*.*.4 | * | Forward A, D |
| *.*.*.4 | * | Forward D |
| 1.2.3.* | * | Forward A, B, C |
| *.2.3.* | * | Forward B, C |
| 1.*.3.* | * | Forward A, C |
| *.*.3.* | * | Forward C |
| 1.2.*.* | * | Forward A, B |
| *.2.*.* | * | Forward B |
| 1.*.*.* | * | Forward A |

Compilation: Takeaways

- Switch classifiers are not compositional – they do not contain enough information
 - Compiler uses an intermediate form which makes use of high-level information
- Classifiers that implement high-level specifications grow in size quickly
 - Compiler uses optimized algorithms and classifier minimization techniques



REACTIVE SPECIALIZATION

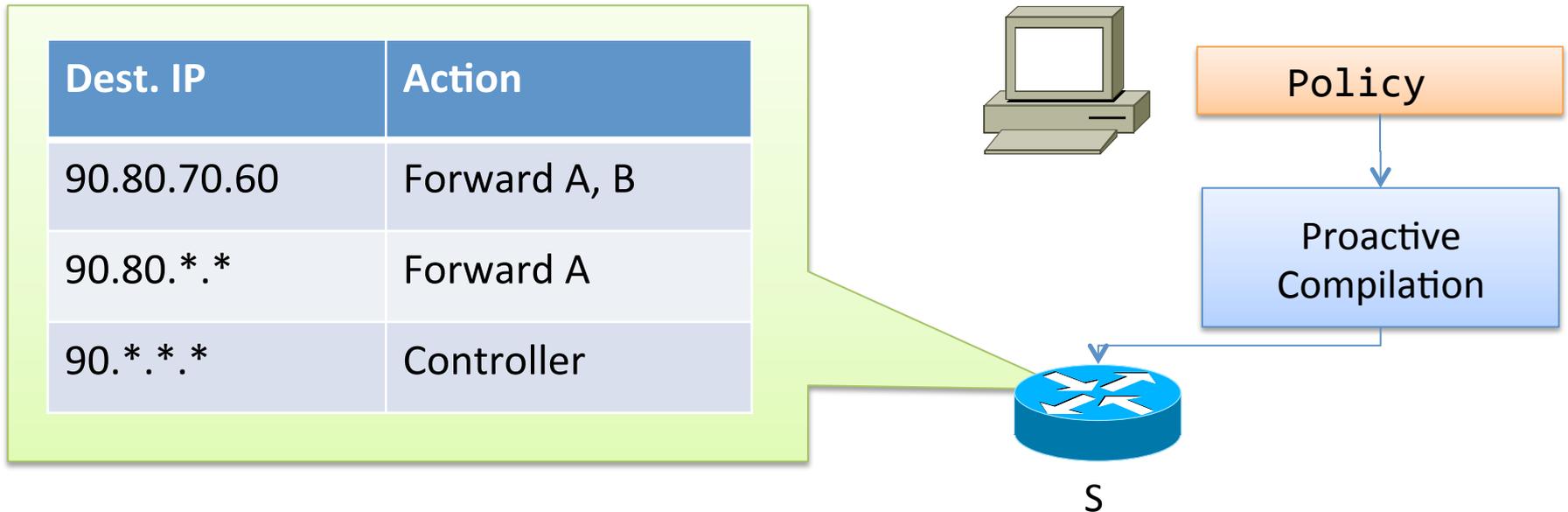
Reactive Specialization

```
DestIP: 90.80.*.* → {A} ∪  
DestIP: 90.*.70.60 → {B}
```

- Can't compile this classifier exactly without using many rules
 - Most switches only support prefix matches
 - Worst case can take billions and billions of rules

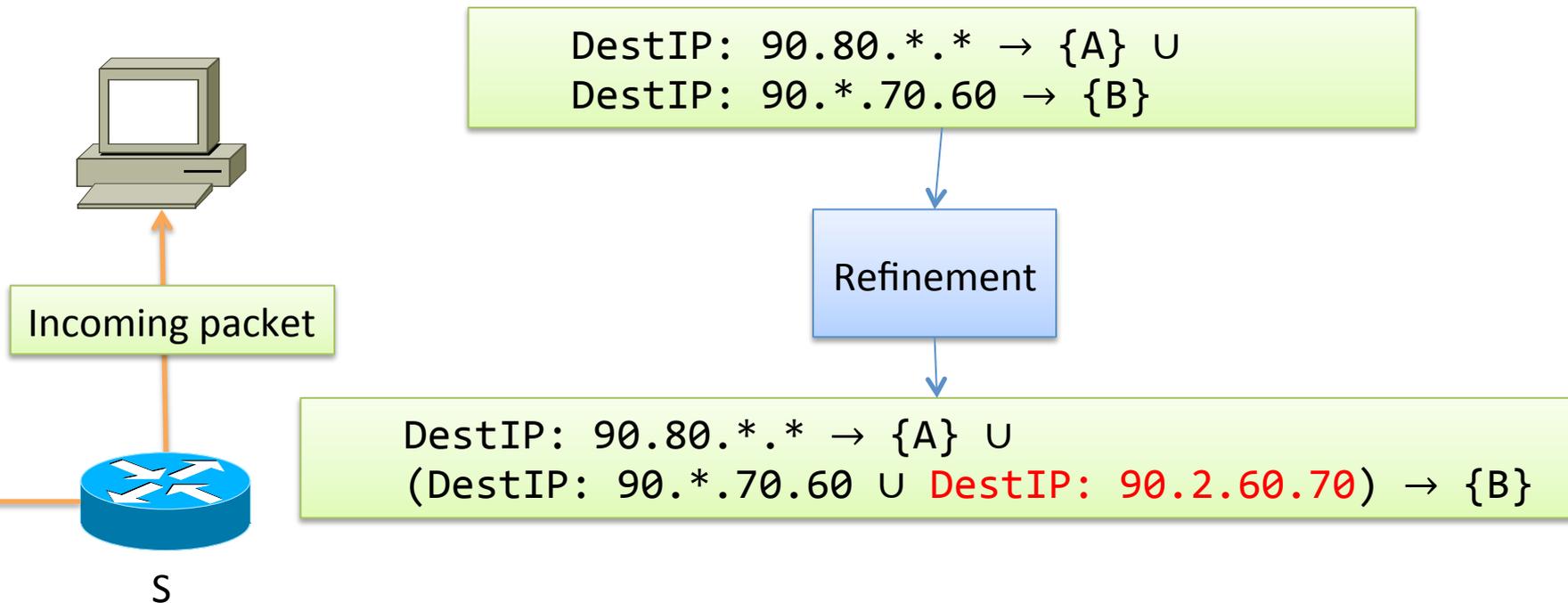
R. Specialization: Base Compilation

- Step 1: install a “base” classifier
 - Proactively handle packets on switch if possible
 - Send other packets to the controller



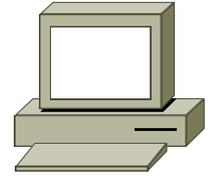
R. Specialization: Refinement

- Packet {DestIP=90.2.60.70, DestPort=80} comes in
- Use packet to generate semantically equivalent, but structurally different policy



R. Specialization: Recompilation

- Compile refined policy
 - Extra structural information handles similar packets to the one we handled



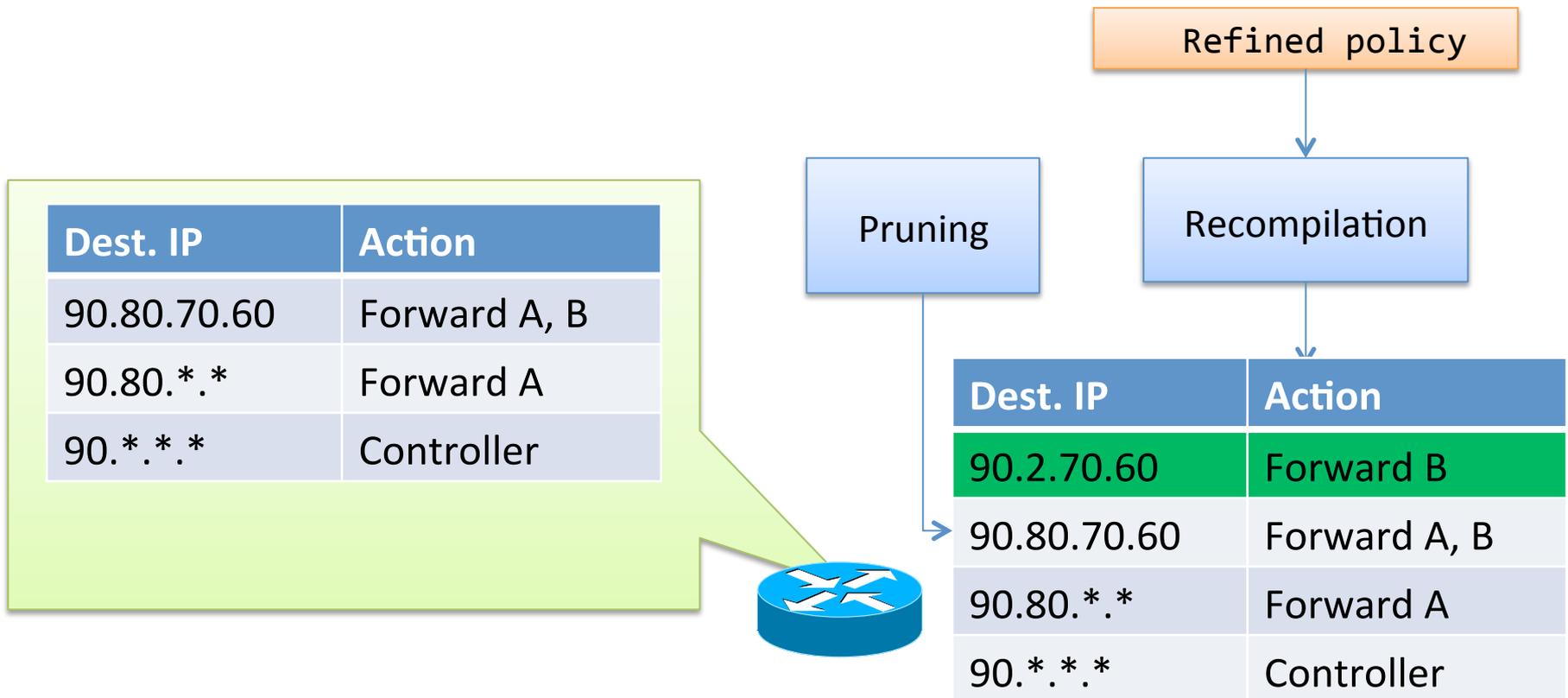
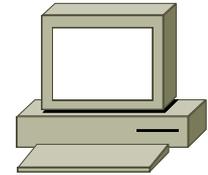
Refined policy

Recompilation

| Dest. IP | Action |
|-------------|--------------|
| 90.2.70.60 | Forward B |
| 90.80.70.60 | Forward A, B |
| 90.80.*.* | Forward A |
| 90.*.*.* | Controller |

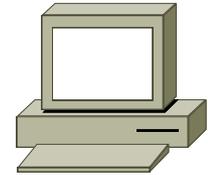
R. Specialization: Pruning

- Remove controller rules and rules that don't match packet we handled
- Take result and place on switches



R. Specialization: Pruning

- Remove controller rules and rules that don't match packet we handled
- Take result and place on switches



Refined policy



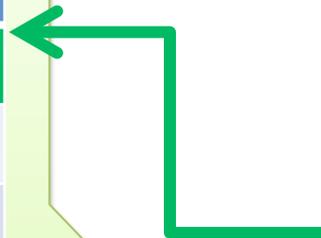
Recompilation



| Dest. IP | Action |
|-------------|--------------|
| 90.2.70.60 | Forward B |
| 90.80.70.60 | Forward A, B |
| 90.80.*.* | Forward A |
| 90.*.*.* | Controller |

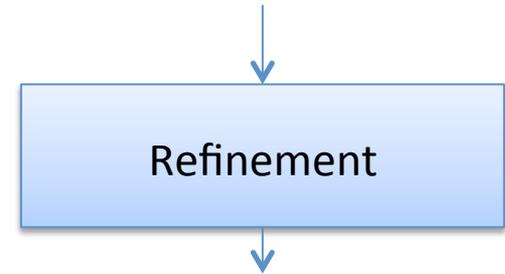
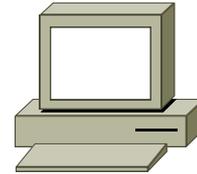


| Dest. IP | Action |
|-------------|--------------|
| 90.2.70.60 | Forward B |
| 90.80.70.60 | Forward A, B |
| 90.80.*.* | Forward A |
| 90.*.*.* | Controller |



R. Specialization: Second Packet

- Consider handling a second packet, {DestIP=90.100.70.60, DestPort=21}



DestIP: 90.80.*.* → {A} ∪
(DestIP: 90.*.70.60 ∪ DestIP: 90.100.70.60) → {B}

| Dest. IP | Action |
|--------------|--------------|
| 90.100.70.60 | Forward B |
| 90.2.70.60 | Forward B |
| 90.80.70.60 | Forward A, B |
| 90.80.*.* | Forward A |
| 90.*.*.* | Controller |



| Dest. IP | Action |
|--------------|--------------|
| 90.100.70.60 | Forward B |
| 90.80.70.60 | Forward A, B |
| 90.80.*.* | Forward A |
| 90.*.*.* | Controller |

THEORETICAL & EMPIRICAL RESULTS

Functional Correctness

- Define idealized machine that handles all traffic at the controller

Theorem (Functional Correctness):

NetCore machines bisimulate idealized machines.

- Statement of correctness more subtle when considering queries

Quiescence

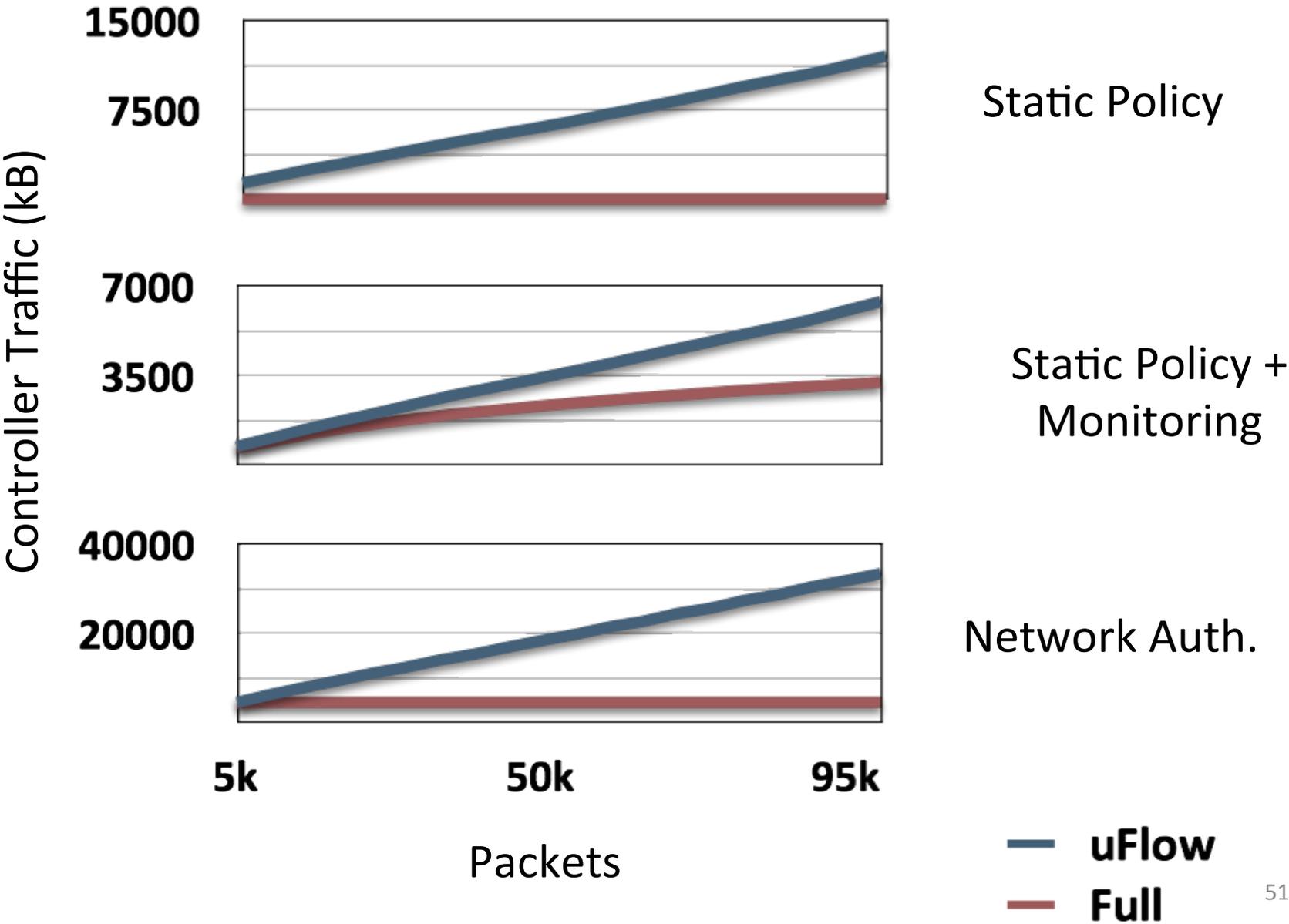
- Recall: Programs are efficient if they handle traffic on the switches, not the controller.

Theorem (Quiescence):

Under reasonable circumstances, NetCore programs eventually handle all traffic on the (fast) switches, instead of on the (slow) controller

- Special case of cost theorem: identify packets sent to controller

Benchmarks: Controller Traffic



Conclusion

- This talk was short; we didn't talk about:
 - Switch heterogeneity & lattice theory
 - Arbitrary Haskell functions inside policies
 - Queries
 - How compilation and specialization actually work
- So:
 - Check out the paper!
 - Talk to Chris, Dave, or Nate!
 - Check out the compiler at frenetic-lang.org!

frenetic >>