

- Limited privilege
- Achieving root access on machine X may require multiple steps
  - Get inside firewall
  - Scan network for vulnerabilities
  - Get user access to machine
  - Get root access to machine
- Question: how does security of whole system depend on parts?

- Must handle large, realistic examples
- Should guarantee properties of attack graphs

– Analysis

- Must enable security analysis by system administrators
- Should support incremental, partial specification

- Handles safety and liveness properties
- *Generates counterexamples*

An **attack graph** is a set of attacks of  $\Pi$ .

[Gerth et al.95].

- $M$  and  $\Phi$  induce languages  $L(M)$  and  $L(\Phi)$ .
2. Compute intersection  $M \cap \sim\Phi$  of Buchi automata.
    - $L(M \cap \sim\Phi) = L(M) \setminus L(\Phi) =$  executions of  $M$  that violate  $\Phi$ .
  3. Derive  $G$  from strongly connected components of intersection automaton [Tarjan72].

IIS buffer overflow:	remotely get root	X
Squid portscan:	port scan	X
LICQ remote-to-user:	gain user privileges remotely	✓
scripting exploit:	gain user privileges remotely	X
local buffer overflow:	locally get root	X

```

     $R(S, T, 80)$ 
intruder effects
     $plvl(T) := \text{root}$ 
network effects
     $\neg w3svc_T$ 
end

```

*Host  $T$  is reachable from  $S$  on port 80*

*Root-level privileges on host  $T$*

*Host  $T$  is not running IIS*

**network preconditions**

$\text{licq}_T$

$R(S, T, 5190)$

**intruder effects**

$\text{plvl}(T) := \text{user}$

**network effects**

$\emptyset$

**end**

*Host  $T$  is running vulnerable LICQ software*

*Host  $T$  is reachable from  $S$  on port 5190*

*User-level privileges on host  $T$*

*No changes to the network component*



```
    plvl(T) := root  
network effects  
     $\emptyset$   
end
```

*Root-level privileges on host *T**

*No changes to the network component*



Attack Graphs

20

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### Solution (Sketch):

1. Reduce MCSA to Minimum Hitting Set (MHS) Problem [JSW02].
2. Reduce MHS to Minimum Set Covering (MSC) Problem [ADG80].
3. Use textbook Greedy Approximation Algorithm to approximate solution [CLR85].

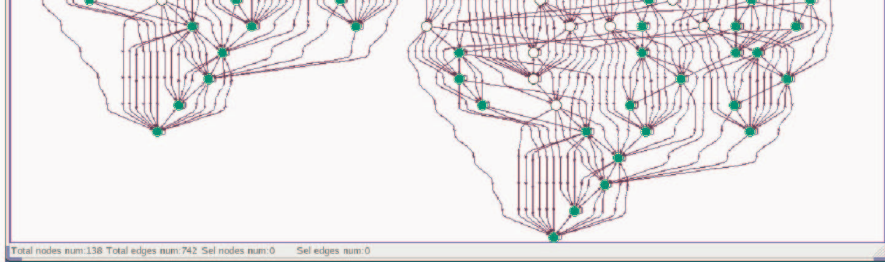
**MCSA** and **MHS** are polynomially-equivalent.

[JSW02b] Jha, Sheyner, Wing, "Two Formal Analyses of Attack Graphs,"  
Computer Security Foundations Workshop, Nova Scotia, June 2002.

Use textbook Greedy Approximation Algorithm for **MSC**  
[CLR85, p. 975.]

deploy to make the system safe: [50+]

**Solution Approach:** Greedy algorithm with provable bounds.  
General case is NP-complete (slightly more complex than minimum cover problem).



```
      <ftp/> <sshd/> </remote>
    </connectivity>
    <cve>
      <CVE_2002_0004/>
      <CVE_2001_1030/>
      <CVE_2001_0439/>
    </cve>
  </host>
```

Attack Graphs

32

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